Lizard

A Linked Data Publishing Platform

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The (a) real world of service provision

What to do about (some of) it

How to do that



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Editor on SPARQL query

A committer on Apache Jena

At Epimorphics Ltd

This work

- > Epimorphics
- ➤ Funding : InnovateUK^{*}
- > Users
 - $\circ~$ For the discussion and encouragement

* Used to be the Technology Strategy Board. UK Department for Business, Innovation & Skills

http://environment.data.gov.uk/

http://landregistry.data.gov.uk/

Customer Requirements

Maximise usage

Publication not application

Data publishing != Database backed web site

• Different traffic patterns

- Expensive queries, less control
- Bot multiplier effect
- "Admin"
 - SLAs: Heartbleed

Problem Statement

- Reacting to events
- Machine administration / SLAs



24x7 Operation

Consistency

Makes the system easier to use

- For users
- For operators

Each query sees an unchanging database

... that did exist; no "bit of this, bit of that"

Clients may conspire!

Apache Jena TDB



➤ Node Table

• Inline values (integers, date/dateTime, ...)

Indexes are covering

- Range scans
- All key, no value
- No "triple table"

SPARQL Execution

```
{ ?x :p 123 . }
```

Convert to Nodelds

Look in POS to get all PO?, assign S to ?x

123 is an inline constant in TDB.

```
{ ?x :p 123 .
?x :q ?v . }
A database join
Index join (Loop+substitution)
Index join (= loop) on
:x1 :q ?v
where :x1 is the value of ?x
```

Index Implementation

TDB uses threaded B+Trees for indexes

8K blocks 100-way B+Tree





Query and Update

Indexes / B+Trees

Blocks

Node table / Objects

 $Key \rightarrow Value Store$

This Does Not Work (very well)



Distribute the storage K->V store Index access on query processor

- Easy to do (pick a KV store of your choice)
 Impedance mismatch

 Too much data moving about
 Little parallelism
 - Bad cold-start



Distribute the indexes

- With modified index access
- Distribute the nodes
- > Comms : Apache Thrift

Clustered Node Table

Node Table

• N replicas; Read R / Write W

e.g. W=N and R =1 =>

Complete copies of node table on each data server

- Can shard
- Replaceable

Requirement: Nodeld for naming

Clustered Indexes

> Indexes

- Can shard by subject
- Replicas of each shard (R=1, W=N)
- Compound access operations

Clustered Indexes



Modified SPARQL Execution

Different unit of index access

o subject + several predicates
 (subj, pred1, pred2, pred3, ...)

Different join algorithms

- Merge join
- Parallel hash join

Configuration 1



Configuration 2





Working prototype

Spin-off : TDB2

New Technology

- Copy-on-write indexes
- New transactional coordinator
- Apache Thrift encoded node table
- Side effect: TDB2
 - Arbitrary scaling transactions
 - Transactional only
 - Space recovery





CommitStrip.com